

Dye - Eigen - Fremlin Reconstruction Theorem

I. Statements, context & Basic facts

Dye '63: G, H two ergodic f.g. on
a standard proba. space (X, μ)
then any $\Psi: G \rightarrow H$ gp. isomorphism
is the conjugation by some p.m.p. bijection.
i.e. $\exists S \in \text{Aut}(X, \mu), \forall T \in G: \Psi(T) = STS^{-1}$

Fremlin '2002: Let \mathcal{A} and \mathcal{B} be two Dedekind
complete Boolean algebras
 G, H two subgroups of $\text{Aut}(\mathcal{A})$ $\text{Aut}(X, \mu)$
and $\text{Aut}(\mathcal{B})$ respectively, both with many
involutions.

if $\Psi: G \rightarrow H$ is a group isom., then

$\exists!$ Boolean isom. $S: \mathcal{A} \rightarrow \mathcal{B}$ s.t. $\forall T \in G:$
 $\Psi(T) = STS^{-1}$.

Definition: $G \leq \text{Aut}(X, [\mu])$ sg. has many involutions,
if $\forall A \subseteq X$ of positive measure \exists non-trivial
involutions $V \in G, \text{supp } V \subseteq A$
" "
 $\{x \in X \mid T(x) \neq x\}$

Proposition: if G is an ergodic f.g. of (X, μ) then G has many involutions.

Proposition: ($\text{MAlg}(X, \mu)$ is Dedekind - complete)

phD manuscript of François. Any family of elements of $\text{MAlg}(X, \mu)$ admits a supremum which is obtained as the limit of an increasing sequence of finite reunions of elements of the family.

Disclaimer: We work "up to measure 0"

by $A \cap B = \emptyset$ I mean that $\mu(A \cap B) = 0$

Lemma: $T \in \text{Aut}(X, [\mu])$, $\emptyset \neq D \subseteq X$, TFAE:

(i) $\text{supp } T \not\subseteq D$

(ii) $T|_{X \setminus D} \neq \text{id}|_{X \setminus D}$

(iii) $\exists \emptyset \neq C \subseteq X \setminus D$ s.t. $C \cap T(C) = \emptyset$

Moreover, $\text{supp } T = \sup \{ D \subseteq X \mid D \cap T(D) = \emptyset \}$

Proposition: $\forall \text{involut}^i \in \text{Aut}(X, [\mu])$, then $\exists A, B \subseteq X$

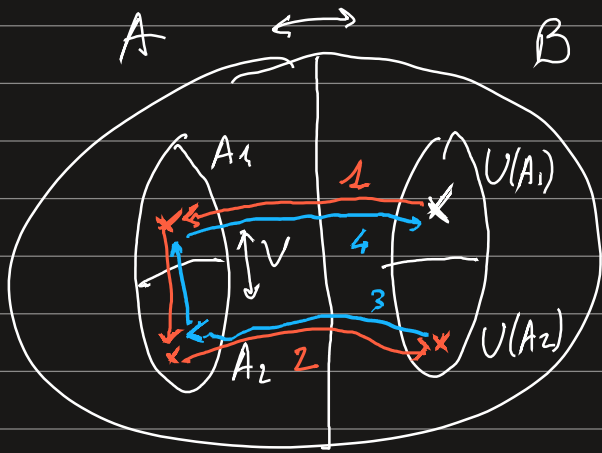
disjoint, such that $\begin{cases} V(A) = B \\ \text{supp } V = A \cup B \end{cases}$
we then say that $V = V_{A,B}$ is an (A, B) -exchanging involution.

II. Linking measure-theoretic properties & g. theoretic properties

Lemma 0: $\mathcal{G} \leq \text{Aut}(X, [\mu])$ w/ many involutions;
 $\forall \emptyset \neq C \subseteq X, \exists T \in \mathcal{G}$ of order exactly 4, s.t. $\text{supp } T \subseteq C$.

Proof:

$U(A, B)$ - exch. (involutions in \mathcal{G} , $A \cup B \subseteq C$)
 non-trivial



$V(A1, A2)$ - exch. Invo.
 w/ $A1 \cup A2 \subseteq A$

$$\text{Supp}(UVU) = U(\text{Supp } V) \subseteq B$$

V and UVU commute
 UVU is an involution

$\Rightarrow T := UV$ has order 4. \square

For the rest of this section, we fix $\mathcal{G} \leq \text{Aut}(X, [\mu])$
 w/ many involutions and V non-trivial involution in \mathcal{G} ,
 w/ $\text{supp } V = A \cup B$
 Define the following:

$$\mathcal{C}_V := C(V) = \{T \in \mathcal{G} \mid TV = VT\}$$

$$\mathcal{D}_V := \left\{ \begin{array}{l} \text{involutions in } \mathcal{C}_V \text{ commuting w/ all their} \\ \mathcal{C}_V\text{-conjugates} \end{array} \right\}$$

$$= \left\{ T \in \mathcal{C}_V \mid T \text{ inv. and } \forall S \in \mathcal{C}_V: TSTS^{-1} = STS^{-1}T \right\}$$

$$\mathcal{E}_V := C(\mathcal{D}_V) = \left\{ T \in \mathcal{G} \mid \forall S \in \mathcal{D}_V: TS = ST \right\}$$

$$\mathcal{F}_V := \text{Sq}(\mathcal{E}_V) = \{T^2 \mid T \in \mathcal{E}_V\}$$

$$\mathcal{G}_V := C(\mathcal{F}_V) = \{T \in \mathcal{G} \mid \forall S \in \mathcal{F}_V: ST = TS\}$$

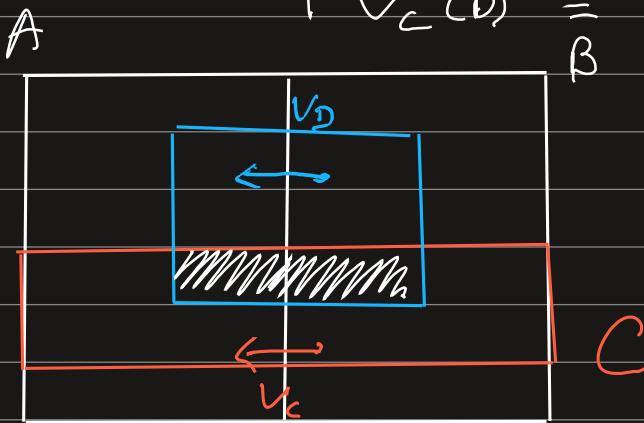
Lemma 6: $\forall T \in \mathcal{G}_V: T(\text{supp } V) = \text{supp } V$

Proof: $T(\text{supp } V) = \text{supp}(TVT^{-1}) = \text{supp}(V) \quad \square$

Definition: $\forall \emptyset \neq C \subseteq X$ s.t. $V(C) = C$

the induced exchanging-involution V_C is defined by

$$\begin{cases} V_C(D) = V(D) & \text{if } D \subseteq C \\ V_C(D) = D & \text{if } D \subseteq X \setminus C \end{cases}$$



Notice that:

$$V_C = V_{(A \cap C), (B \cap C)}$$

$$V_C V_D = V_{C \triangle D} = V_D V_C$$

Lemma 7: (1) $\forall T \in \mathcal{D}_V: \text{supp } T \subseteq \text{supp } V$

(2) $\forall \emptyset \neq C \subseteq X$; s.t. $V(C) = C$
 $V_C \in \mathcal{D}_V$.

Proof: (1) By contraposition.

Let $T \in \mathcal{G}_V$ s.t. $\text{supp } T \not\subseteq \text{supp } V$



$\emptyset \neq C \subseteq \text{supp } T$ s.t.

$$C \cap T(C) = \emptyset$$

Lemma 0 $\Rightarrow \exists S \in G$ of order 4, $\text{supp } S \subseteq C$

• $SV = VS \Rightarrow S \in G_V$

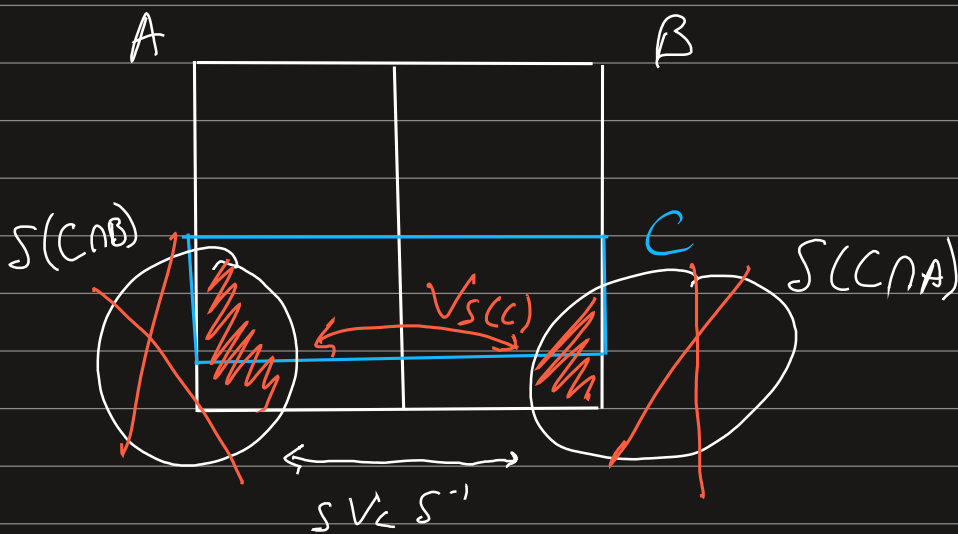
• $S \neq S^{-1}$ so $\exists \emptyset \neq D \subseteq C, S(D) \neq S^{-1}(D)$

we also have, $C \cap T(D) = \emptyset = C \cap TS^{-1}(D)$

$\rightarrow ST \cancel{S^{-1}} T(D) = S \cancel{T}^{-1}(D) \neq S^{-1}(D) = T \cancel{TS^{-1}}(D)$

so (D non-trivial) T and STS^{-1} do not commute.
 $T \notin \mathcal{D}_V$.

(2) $\forall S \in \text{Aut}(X, [\mu]) : SV_C S^{-1} = (SVS^{-1})_{S(C \cap A), S(C \cap B)}$



In particular, for $S = V, VV_C V^{-1} = V_C \Rightarrow V_C \in G_V$.

Now for any $S \in G_V, \underbrace{VS(C)}_1 = SV(C) = \underbrace{S(C)}_2$

$\underbrace{SV_C S^{-1}}_1 = (SVS^{-1})_{S(C \cap A), S(C \cap B)} = V_{S(C \cap A), S(C \cap B)} = V_{S(C)}$

$\underbrace{V_{S(C)}}_1 V_C = V_C \underbrace{V_{S(C)}}_2$ i.e. $V_C \in \mathcal{D}_V$. \square

Lemma 8:

(1) $G_V \subseteq G$

(2) $\forall T \in G_V : \forall \emptyset \neq C \subseteq \text{supp } V : T(C) \subseteq C \cup V(C)$

$C \cup V(C)$

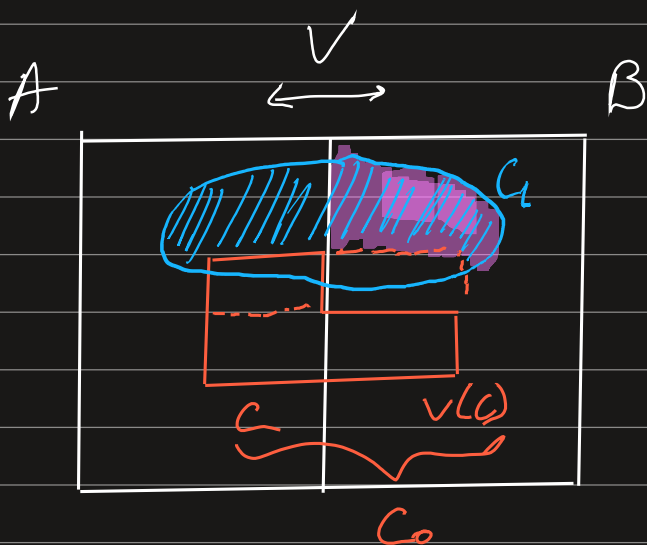
$$(3) \forall T \in \mathcal{E}_V : \forall \emptyset \neq C \subseteq \text{Supp } V : T^2(C) = C$$

(4) if $T \in \mathcal{G}$ is s.t. $\text{supp } T \cap \text{supp } V = \emptyset$,
then $T \in \mathcal{E}_V$.

Proof: (1) $V \in \mathcal{D}_V$. ✓

(2) Assume the contrary: $\emptyset \neq C \subseteq \text{Supp } V$ s.t.

$$T(C) \not\subseteq C \cup V(C)$$



$$V(C_0) = C_0$$

$$\emptyset \neq C_1 := T(C_0) \setminus C_0$$

(1) & Lemma 6

$$\Rightarrow C_1 \subseteq T(\overbrace{A \cup B}^{\text{supp } V}) = A \cup B$$

$$VT(C_0) = TV(C_0) = T(C_0)$$

$$V(C_1) = VT(C_0) \setminus V(C_0) = T(C_0) \setminus C_0 = C_1$$

Define $D := B \cap C_1$ notice 2 things:

⊙ $C_1 = V(D) \cup D$

⊙ $C_1 \cap T(D) = \emptyset : C_1 \subseteq T(C_0) \Rightarrow C_1 \cap T(D) \subseteq T$

Therefore, $V_{C_1} T(D) = T(D) \neq TV_{C_1}(D)$

But $V_{C_1} \in \mathcal{D}_V$ ✓

$$(3) \quad \text{supp } T = \sup \{ D \subseteq X \mid D \cap T(D) = \emptyset \}$$

$$D \subseteq C \cap \text{supp } T = \sup \{ D \subseteq C \mid D \cap T(D) = \emptyset \}$$

$\xrightarrow{\text{s.t. } D \cap T(D) = \emptyset} C \cap \text{supp } T \subseteq C \subseteq \text{supp } V$ so by applying (2) to $C \cap \text{supp } T$ we get

$$\rightarrow T(D \cap \text{supp } T) \subseteq \underbrace{(D \cap \text{supp } T)} \cup \underbrace{V(D \cap \text{supp } T)}$$

in fact $T(D) \subseteq V(D)$

$$\rightarrow T(C \cap \text{supp } T) \subseteq V(C \cap \text{supp } T)$$

by supremum.

We want $T^2(C) = C$

① $VT = TV$ so

$$\uparrow T^2(C \cap \text{supp } T) \subseteq TV(C \cap \text{supp } T)$$

$$\parallel VT(C \cap \text{supp } T) \subseteq V^2(C \cap \text{supp } T)$$

② $\text{supp } V = T^2(\text{supp } V) \quad T \in \mathcal{E}_V \subseteq \mathcal{E}_V$

$$\text{supp } V \setminus T^2(C) = T^2(\text{supp } V \setminus C) \subseteq \text{supp } V \setminus C$$

$$\rightarrow T^2(C) \supseteq C$$

(4) immediate from Lemma 2. (1) □

Lemma 7:

$$S \in \mathcal{F}_V \Rightarrow \text{supp } S \cap \text{supp } V = \emptyset$$

Lemma 8 (3)

Moreover, $\forall \emptyset \neq C \subseteq X \mid \text{supp } V, \exists$
(involution non-trivial) $S \in \mathcal{F}_V$, w/ $\text{supp } S \subseteq C$

Proof: Lemma 6 $\Rightarrow \exists T \in \mathcal{G}, \text{supp } T \subseteq C$

Lemma 5 (4) $\Rightarrow T \in \mathcal{E}_V$. T^2 is as wanted. ~~QED~~

Lemma 7: $(\mathcal{G}_V := \{T \in \mathcal{G} \mid \text{supp } T \subseteq \text{supp } V\})$

Proof: \supseteq easy. fix $T \in \mathcal{G}, \text{supp } T \subseteq \text{supp } V$

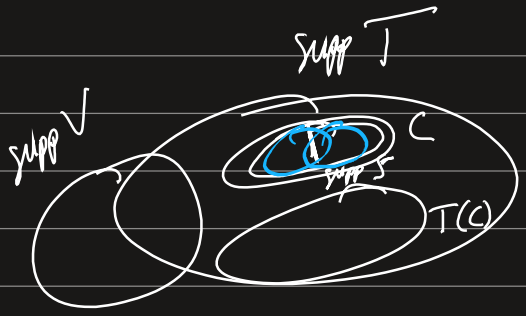
Lemma 7 $\Rightarrow \forall S \in \mathcal{F}_V, \text{supp } S \cap \text{supp } V = \emptyset$

so $TS = ST$ \checkmark

\subseteq By contraposition: fix $T \in \mathcal{G}, \text{supp } T \not\subseteq \text{supp } V$

fix $\emptyset \neq C \subseteq \text{supp } T \setminus \text{supp } V$ s.t. $C \cap T(C) = \emptyset$

Lemma 7 $\Rightarrow \exists S$ non-trivial inv $\in \mathcal{F}_V$
 $\text{supp } S \subseteq C$



for any $D \subseteq \text{supp } S$ s.t.

$$D \neq S(D)$$

$$TS(D) \neq T(D) = \underline{ST(D)}$$

$\Rightarrow T \notin \mathcal{G}_V$ ~~QED~~

III. Constructing the conjugation

$$\psi : \mathcal{G} \rightarrow \mathbb{H}$$

$$\psi(\mathcal{L}_v) = \mathcal{L}_{\psi(v)}$$

$$\forall T \in \mathbb{G} : \text{supp } T \subseteq \text{supp } V \stackrel{(*)}{\iff} \text{supp } \psi(T) \subseteq \text{supp } \psi(V)$$

$$S, S^* \in (\text{Aut}(X, (\mu)))$$

$$\begin{cases} S(C) = \sup \{ \text{supp } \psi(V) \mid V \in \mathbb{G} \text{ invo, s.t. } \text{supp } V \subseteq C \} \\ S^*(C) = \sup \{ \text{supp } \psi^{-1}(V) \mid V \in \mathbb{H} \text{ invo, s.t. } \text{supp } V \subseteq C \} \end{cases}$$

It remains to show :

- easy $\left\{ \begin{array}{l} \bullet S, S^* \text{ are well-defined.} \\ \bullet S^* = S^{-1} \end{array} \right.$
- easy $\rightarrow \bullet S \text{ is the desired conjugation.}$
- $\bullet \text{ unique in that regards.}$

Lemma:

$$\text{from } (*) \quad \left(\begin{array}{l} \forall A \subseteq X \\ \forall V \text{ invo } \in \mathbb{G} \end{array} \right) : \text{supp } \psi(V) \subseteq S(A) \iff \text{supp } V \subseteq A$$